


The number of Android malware has been increasing dramatically in recent years. Android malware can violate users' security, privacy and damage their economic situation. Study of new malware will allow us to better understand the threat and design effective anti-malware strategies. In this paper, we introduce a new type of malware exploiting Android's accessibility framework and describe a condition which allows malicious payloads to usurp control of the screen, steal user credentials and compromise user privacy and security. We implement a proof of concept malware to demonstrate such vulnerabilities and present experimental findings on the success rates of this attack. We show that 100% of application launches can be detected using this malware, and 100% of the time a malicious Activity can gain control of the screen. Our major contribution is two-fold. First, we are the first to discover the category of new Android malware manipulating Android's accessibility framework. Second, our study finds new types of attacks and complements the categorization of Android malware by Zhou and Jiang [32]. This prompts the community to re-think categorization of malware for categorizing existing attacks as well as predicting new attacks.

Android, Malware, Attack

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1. Introduction

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H Q R U P R X W O K S D W Z R D Z K L O R E L C H Y L F I N F O R M A T I O N
K D Y E H F R P X E L T X W R I R O D L O L I B O D U F K
- X Q L S H W Z R U N O M S R W W H E B L T K H U R D W W P U
K D G L V F R R M H H W K R X F D Q Z D D P S D H V
S H U F I H Q F W R D M H : L W K S H U F R I H Q W E L O H
P D O Z D D H S Q U R D C Z D D B O D Q Q D W H J R U L L Q J
W K H P D O Z D U H P S R U W W D H S V Z S U G L G M Z L Q J
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interface (API) methods to provide customized accessibility services in their own applications. However, the accessibility service has access to critical sensitive information, including information about applications that are currently running and account information. Attackers could utilize such a vulnerability to conduct various types of attacks.

To prove the concept, we develop a malicious application which exploits this vulnerability. The installation, activation, and the payload of our malicious application are described as follows. First, to install our malicious application onto user devices, our malicious application may appear as a legitimate accessibility service application and provide some accessibility functionality. The installation of our malicious application requests the `BIND_ACCESSIBILITY_SERVICE` permission. Of course, other permissions are required if the malicious payload requires such permissions. The malicious application is triggered once an `AccessibilityEvent` object is dispatched. There are

the Android

Accessibility

Framework

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types in section 2. Our malicious application supports various malicious payloads. In this paper, we implement one payload which launches a masquerade attack to collect user's email login credentials.

Our major contributions are summarized below:

- We are the first to identify malware leveraging Android's accessibility framework. The impact of this attack is severe, given that it can be used to activate various malicious payloads. For example, the payload can be various masquerade attacks, emulating Email, Facebook and other popular apps to steal user credentials and cause other damage. We test our proof-of-concept malware against multiple anti-malware applications and none are able to detect it.
- We identify a potential logic error in the way the Android user interface framework manages the order of application launches. When two applications are launched nearly simultaneously, both launch requests consider the first request to be processed as the next application to be launched, preventing the second launch from occurring. This can lead to a variety of attacks, such as denial of service and masquerade attacks, as we discuss in Section 3.3.
- Our study complements the categorization of Android malware by Zhou and Jiang [33], who present a systematic characterization of existing Android malware by their strategies of installation, activation, payload and permission use. We identify new events for activation and our attack shows that the payload can be other attacks, such as the masquerade attack demonstrated in this paper.

- We also present a novel Android application framework, called Android-App-Builder, which can be used to create Android applications that can be used to test the effectiveness of anti-malware applications.

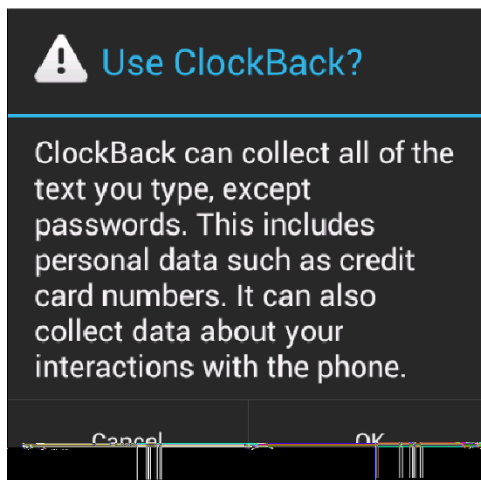


Figure 1. Accessibility Service Warning Dialog

LLED,TYPE_VIEW_SELECTED,TYPE_VIEW_TEXT_CHANGED,TYPE_VIEW_TEXT_SELECTION_CHANGED,TYPE_VIEW_TEXT_TRAVERSED_AT_MOVEMENT_GRANULARITY,TYPE_WINDOW_CONTENT_CHANGED,TYPE_WINDOW_STATE_CHANGED [10].

Additionally, as of Android 1.6, developers may create custom accessibility services by extending the `AccessibilityService` class [11]. Descendants of `AccessibilityService` must override the `onAccessibilityEvent` method, which gets called each time an Accessibility Event occurs and is passed the populated `AccessibilityEvent` object. A custom `AccessibilityService` may then use the information contained in the `AccessibilityEvent` or optionally query the window content for the contextual data needed to perform its function. Naturally, the receipt of an event and its associated data and the ability to query window content present a security risk. For example, every time a user inputs text into an `EditText` widget, a custom `AccessibilityService` will receive an `AccessibilityEvent` with type `TYPE_VIEW_TEXT_CHANGED` that contains the text that the user input. To mitigate this risk, the Android system requires that `AccessibilityServices` be enabled manually by the user and displays a dialog window alerting the user to the risk.

The alert in Figure 1 states that the service can collect all of the text a user types, except for passwords. Passwords are prevented from being collected if the `android:password` XML attribute of the `EditText` object is true, causing the typed characters to display as asterisks and the `AccessibilityEvent` not to be populated with the `EditText` content. We describe a technique in Section 3 that does allow passwords to be collected.

2.2. Masquerade Attack

We formally define masquerade attacks as attacks impersonating apps such as those found on Google Play and other Android markets. Masquerade attacks can use the same names, icons and/or interfaces as those legitimate apps to masquerade. The “or” is used here since it can be enough for a fake app using the same name or icon to attract users to download and commit damage. Of course malware with the same name, icon and interface and functionality similar to the legitimate app is more attractive and more likely to be downloaded and installed.

Many fake applications masquerade as popular games, which include Angry Birds Space [8], Fruit Ninja, Temple Run and Talking Tom Cat. Other popular applications such as Opera Mini [7] are also frequently utilized to launch masquerade attacks. One malware example is Fake Netflix [4] which disguises as Netflix. To analyze its behaviour, we use `apktool` [2] and JD [5] to disassemble the apk file. We find that when the application is launched, a login interface will show and prompt the user to input account information. After the user input account name and password, the application creates a new `HttpClient` and `Post Header`, and send these account information to a remote site `http://erofolio.no/ip/biz/login.php` via an HTTP POST request. After that, a dialog pops up and presents a message “Your Android TV is not supported”. When the user pushes the button on the dialog, the application creates an intent to uninstall itself. Another example is Fake Media Player [1]; we also decompile its apk file to analyze its behaviour. From its `AndroidManifest.xml` file, we find that this application requests a permission to send SMS messages during installation. From its `MoviePlayer.class` and `DataHelper.class` files, we find that the application performs several steps. First, it creates a database. Then, it sends out SMS messages. After that, this application inserts an entry into the database with value ‘was’. The next time the application is launched, it will check if such entry exists. If the entry does not exist, it will send out SMS messages.

Repackaged applications that contain malicious functions are more difficult for users to notice, because they perform all functions that the legitimate applications perform. When a user is using these fake applications, the malicious function is running in the background. One example is a group of wallpaper-related applications. These applications upload user information to a remote server or run searching operations in the background to improve the ranking of certain websites [6].

3. Malware Exploiting Accessibility Service

In this section, we first give an overview of the malware that exploits the Android accessibility framework. We

then address major challenges for such malware to work, including detection of the launch of a victim app and race condition between the victim app and malware.

3.1. Overview

We now introduce the novel malware's installation, activation, malicious payloads and permission uses.

I : The new Android malware can provide regular accessibility service as it claims and conduct attacks silently. Therefore, impaired users and users who prefer large font text may be interested in such malware and install it onto their device. It is also reasonable to assume that the malicious application could be marketed under a category other than accessibility services. For instance, malware authors could market a "driving mode" app which leverages the system's accessibility features in order to provide better hands-free operation while driving an automobile. This installation strategy is *installation.others.3rdgroup* in [33], referring to apps that intentionally include malicious functionality. For brevity, we denote *installation.others.3rdgroup* as *trojan*, "a program made to appear benign that serves some malicious purpose" according to the taxonomy in [27], although this definition of trojan may be controversial.

D : During installation, the malware requests the `BIND_ACCESSIBILITY_SERVICE` permission. After installation, users must enable the AccessibilityService in Android's Accessibility Settings menu. Since a legitimate accessibility service also requests such permission and requires enabling, users may not suspect the motivation of our malware. Of course, other permissions are required if the malicious payload requires them. However, Felt *et al.* [3] show that only 17 percent of Android users actually pay attention to application permissions at the time an app is installed. Furthermore, only 56.7 percent of participants in the study claimed they had canceled an app installation because of issues with its permissions. Given these startling statistics, the installation of a malicious app requiring the `BIND_ACCESSIBILITY_SERVICE` permission and its required payload permissions could realistically be performed by typical Android users.

A : After the installation, our malware can derive a list of all installed applications in that device. This can be achieved via many sources, such as the Package Manager. Based on which applications are installed, our malicious application could download various payloads and use them to launch different attacks. 98102547type54254pd1c546006(546a)19pkhiv c546.Orig97(a)7.ackagy

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- How can the malware display itself to the user while the victim app is hidden in the background? When the user touches an app, this app will be launched. Which app, our impostor or the victim app, will be displayed? How can our impostor win the race condition?

We address these two issues below.

3.2. Detecting Application Launch

A crucial piece of our masquerade attack is the ability to detect the launch of a victim application. There is no public API to allow a user application to be notified when another application is launched. An application could poll the `ActivityManager` for changes in the running task list, but this solution could impact CPU and battery performance, and some delay between launch and the detection could occur. Another technique used by malware authors involves using the [READ_LOGS](#)

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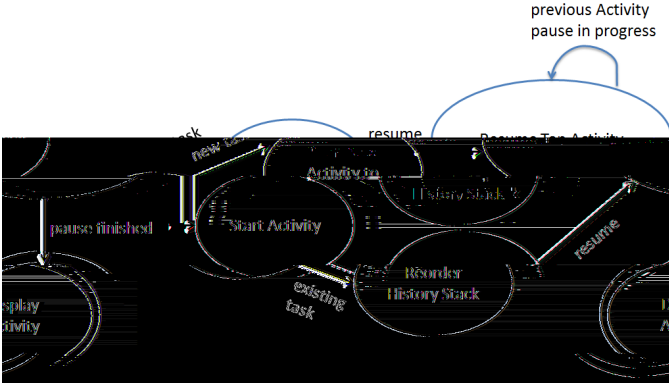


Figure 2. Activity Launch State Transition Diagram

record is moved to the top of the stack. The component which requests the Activity to be started can supply additional flags to the `startActivity` Intent that may alter this behavior, but we focus on the default behavior. Once the stack manipulation is complete, the system attempts to display/resume the top Activity. However, to comply with Android’s Activity lifecycle model, the current Activity must first be paused. This ensures that the current Activity has a chance to save its state before being put into the background. To protect consistency, no Activity is allowed to be displayed until the current Activity has finished pausing. When the pause is complete, the system will attempt to display/resume the Activity at the top of the history stack.

Our malware exploits this logic by launching the malicious Activity as soon as the victim Activity launch has been detected. In the case where both the malicious and victim Activities belong to tasks on the history stack, the malicious Activity will always be displayed over the victim. This is because the launch detection reacts to the `AccessibilityEvent` that is dispatched when the user clicks the Email application icon, which occurs before the Launcher application dispatches its start request. Therefore, the malicious request is received at the `ActivityManager` before the victim request, thus its history record is moved to the top of the stack. Because the two requests are received nearly simultaneously, the victim request is skipped due to the logic in the `ActivityStack` class. This is the fundamental property which makes this attack work when both the victim and malicious Activities are already on the history stack. However, when neither the malicious or victim Activities are on the history stack, the victim Activity is pushed onto the stack after the malicious Activity, making it next in line to be displayed. Figure 3 shows a timeline of when these events occur, illustrating the state of the history stack for each case over time.

Table 2. List of Device Screen Statuses

	With flash	Without flash
Victim interface shows up	Ω_1	Ω_2
Fake interface shows up	Ω_4	Ω_3

In Section 3.4 we show that by adding some delay to the malicious Activity launch, we can increase the chance that the malicious Activity will be pushed onto the stack after the victim.

3.4. Optimal Delay

There exists a source of contention for our malicious Activity. An attacker wants the malicious Activity to be displayed instead of the victim Activity without any suspicious screen flash, flicker, or transition animation that may alert the user to the presence of malware. To achieve this goal, the timing of launching malicious Activity should be carefully adjusted so that the malicious Activity is processed soon after the victim Activity. Therefore, the problem is how to derive an optimal delay for the malicious Activity. We present our analysis below.

We find that different delay of the malicious Activity produces four different statuses of the device screen. Before introduce the four statuses, please note that when the malicious Activity is processed, a fake interface is created and displayed. Please also note that when the victim Activity is processed, a victim interface is created and displayed. Depending on the timing of processing each activity, there are four scenarios. (I) The malicious Activity is processed before the victim Activity. The fake interface will be displayed first, and then replaced by the victim interface. In this scenario, we can observe the victim interface showing up with a flash, and the status of device screen is defined to be Ω_1 . (II) The malicious Activity is processed before the victim Activity, but these two activities are processed nearly simultaneously. In this scenario, only the victim interface is displayed without a flash. The status of device screen is defined to be Ω_2 . (III) The malicious Activity is processed after the victim Activity, but these two activities are processed nearly simultaneously. In this scenario, only the fake interface is displayed without a flash. The status of device screen is defined to be Ω_3 . (IV) The malicious Activity is processed after the victim Activity. The victim interface will be displayed first, and then replaced by the fake interface. In this scenario, we can observe the fake interface showing up with a flash, and the status of device screen is defined to be Ω_4 . These statuses are listed in Table 2.

It is obvious that an attacker will want the device screen status to be Ω_3 . Therefore, we need to derive an optimal delay time for the malicious Activity.

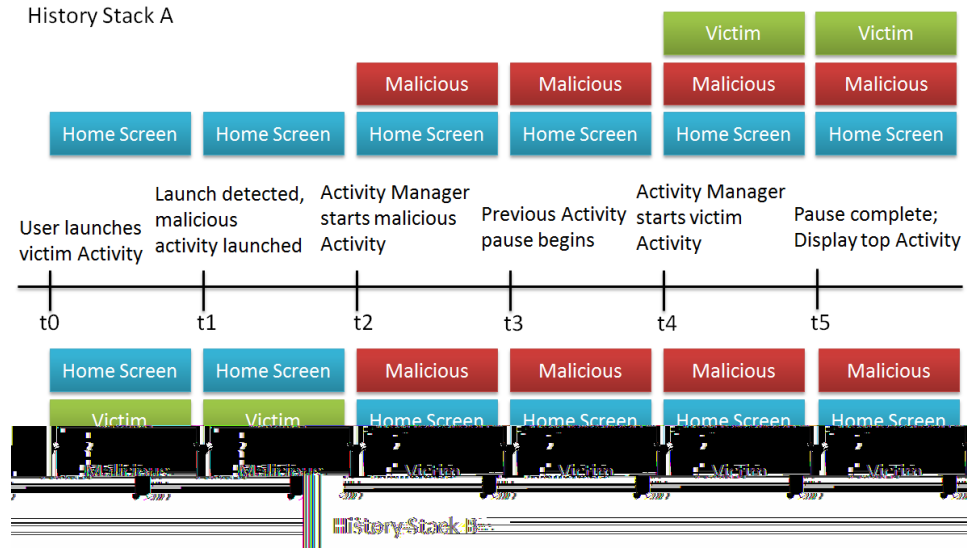


Figure 3. Activity History Stack Timeline

However, for every victim app, the optimal delay time is different. To derive the optimal delay for a specific app, we can enumerate possible delay d to start the malicious Activity with delay d . The delay values producing the highest probability that status Ω_3 occur are the optimal ones.

First, we select a set of different delays, which are denoted as $d_1, d_2, d_3, \dots, d_n$. Next, we select delay d_1 and start the malicious Activity with delay d_1 . We repeatedly launch the victim app a number of times and record the number of occurrences of each status during these tests. We then select the next delay and repeat the second step until all reasonable delays are tested. Finally, we select a range of delay times within which the number of occurrences of status Ω_3 becomes maximal, and correspondingly, the delays within this range are the optimal delays. We evaluate this strategy in Section 4.

Using the strategy above, a malware author can derive an optimal delay for every victim app running on different device models with different Android versions, and keep a record of these optimal delays in a remote server which provides all the malicious payloads we developed. When our malicious service tries to download a malicious payload from that server, the malicious service also sends a request to get an optimal delay to trigger that payload. Therefore, every time a victim app is launched, a fake interface is always shown without a flash.

4. Evaluation

To prove malware exploiting Android's accessibility framework is a new breed, we test our proof-of-concept implementation against four free Android mobile security applications from the Google Play Market:

AVG AntiVirus, Norton Mobile Security, Lookout Mobile Security, and Trend Micro Mobile Security. None of the four detected any suspicious behavior or raised a flag during virus scanning while the AccessibilityService was enabled.

In the following, we provide the experimental results for detecting application launch and winning the race condition. The malicious AccessibilityService was tested on an Android emulator running Android 4.2, an HTC Nexus One running Android 2.3.6, and a Samsung Galaxy S2 running CyanogenMod 9.

4.1. Detecting Application Launch

In the Android versions we tested, application launch detection can be done for any application that defines a Launcher Activity. This is true because the objects that represent shortcut icons on the home screen and the app drawer are descendants of the TextView class, and do not override the behavior of the onPopulateAccessibilityEvent method. The TextView's onPopulateAccessibilityEvent method adds the character string contained in its text field to the text list of character strings in the AccessibilityEvent. The BubbleTextView class used to represent shortcut icons will always store the title of the shortcut that is displayed in the home screen/app drawer in its text field. Therefore, the launch detection simply must match this title in order to detect application launch. Application shortcuts that belong to the hotseat, the horizontal space at the bottom of the default home screen that stays "docked" when navigating to alternate home screens, do not display a title but all shortcut icons in the application drawer do. Although the hotseat applications do not display a title, the AccessibilityEvent that is dispatched on

click does contain the title, however this is not the case in CyanogenMod 9.

We tested the launch detection capability on the HTC Nexus One for the following six Android applications: Messaging, Email, LinkedIn, Facebook, Bank of America, and Browser. For each application, a shortcut icon was created on the home screen. The launch detection was successful for all six applications.

4.2. Winning Race Condition

During the launch detection testing, we noticed that the malicious Activity was not displayed instead of the victim Activity 100% of the time, especially when the victim application was being launched for the first time since system boot. To test this, we performed two separate experiments. In the first experiment, we ensured the application we were launching was not running by pressing the *Force Stop* and *Clear Data* buttons under the corresponding *Settings -> Manage Applications -> All menu* for that application. These two operations effectively force the application to be reloaded from its initial state, as if the system had just booted. We then returned to the home screen, launched the application normally, and recorded which Activity, malicious or victim, was displayed. In the

Table 4. Status with Reloading (HTC Nexus One running Android 2.3.6)

Delay (ms)	Ω_1	Ω_2	Ω_3	Ω_4
0	0%	100%	0%	0%
1 - 740	0%	0%	100%	0%
750	0%	0%	90%	10%
760	0%	0%	70%	30%
770	0%	0%	50%	50%
780	0%	0%	30%	70%
790	0%	0%	30%	70%
800	0%	0%	20%	80%
810	0%	0%	20%	80%
820	0%	0%	20%	80%
830	0%	0%	10%	90%
840	0%	0%	0%	100%

Table 5. Status without Reloading (HTC Nexus One running Android 2.3.6)

Delay (ms)	Ω_1	Ω_2	Ω_3	Ω_4
0	0%	20%	80%	0%
1-60	0%	0%	100%	0%
70	0%	0%	90%	10%
80	0%	0%	100%	0%
85	0%	0%	80%	20%
90	0%	0%	90%	10%
95	0%	0%	80%	20%
100	0%	0%	80%	20%
105	0%	0%	90%	10%
110	0%	0%	50%	50%
115	0%	0%	10%	90%
120	0%	0%	20%	80%
125	0%	0%	10%	90%
130	0%	0%	20%	80%
135 - 145	0%	0%	0%	100%

All menu for the Browser app. These two operations effectively force the application to be reloaded from its initial state when it is launched. (III) We launch the Browser app from home screen, and count the number of screen statuses we observe. (IV) For each delay, we repeat step (II) and (III) many times and calculate the rates of occurrences of those four statuses defined in Table 2. We present our experimental results in Table 4.

From Table 4, we make the following observations. (I) When delay is 0ms, the victim interface shows up without a flash (Status Ω_2) 100% of the time. (II) When delay increases from 1ms to 740ms, the fake interface shows up without a flash (Status Ω_3) 100% of the time. (III) When delay increases from 750ms to 840ms, the rate of (Status Ω_3) decreases to 0%, while the rate of fake interface showing up with a flash (Status Ω_4)

Table 6. Status with Reloading (Samsung S4 running Android 4.4.4)

Delay (ms)	Ω_1	Ω_2	Ω_3	Ω_4
0	0%	100%	0%	0%
1-10	0%	97.5%	0.025%	0%
15	0%	50%	50%	0%
20	0%	10%	90%	0%
25-115	0%	0%	100%	0%
120	0%	0%	80%	20%
140	0%	0%	20%	80%
150-500	0%	0%	0%	100%

increases to 100%. Therefore, the optimal range of delay is between 1ms and 740ms.

In the second group of tests, we test the Browser app without reloading. The procedure for our test is as follows. (I) We set the delay of our malicious Activity to different values. (II) We launch the Browser app from the home screen, and record the status of the screen we observe. (III) For every setting of delay, we repeat step (II) a number of times, and calculate the rates of those four statuses defined in Table 2. We present our experimental results in Table 5.

From Table 5, we make the following observations. (I) When delay is 0ms, 20% of the time the victim interface shows up without a flash (Status Ω_2). 80% of the time the fake interface shows up without a flash (Status Ω_3). (II) When delay increases from 1ms to 60ms, the fake interface shows up 100% of the time without a flash. (III) When delay increases from 70ms to 105ms, the fake interface shows up 80% to 90% of the time without a flash, and the fake interface shows up 10% to 20% of the time with a flash (Status Ω_4). (IV) When delay increases from 110ms to 145ms, the probability of Status Ω_3 decreases to 0%, while the probability of Status Ω_4 increases to 100%. Therefore, we derive a range of optimal delay, which is between 1ms and 60ms.

In the first group of experiments, the range of optimal delay is between 1ms and 740ms. In the second group of experiments, the range of optimal delay is between 1ms and 60ms. Combining the experimental results from these two groups, the optimal delay for the malicious Activity is between 1ms and 60ms.

To illustrate the difference in performance between older and current devices and Android versions, we performed a set of experiments on a Samsung Galaxy S4 running Android 4.4.4. The optimal delay did change between the older devices and the more modern S4, likely due to the S4's faster hardware and newer OS. Table 6 summarizes the characteristics of the delay on the S4 when clicking the Browser application with reloading, while Table 7 summarizes the results without reloading. There is a significant difference between the S4 delays and the older devices in Tables 4 and 5,

Table 7. Status without Reloading (Samsung Android 4.4.4)

Delay (ms)	Ω_1	Ω_2	Ω_3	Ω_4
0	0%	100%	0%	0%
1	0%	90%	10%	0%
10	0%	100%	0%	0%
20	0%	20%	80%	0%
30-200	0%	0%	100%	0%
300	0%	0%	40%	60%
400	0%	0%	10%	90%
500	0%	0%	0%	100%

particularly for small delay values. Considering the results from all four tables, the optimal delay for the malicious Activity is between 30ms and 60ms. To avoid suspicion in the event that the improved performance of an OS upgrade alters the optimal delay, a malware author could restrict the payload activation to only those versions which have a known optimal delay.

5. Discussion

In this section, we discuss various ways that the malware discussed in this paper can be enhanced and become much more sophisticated. Given the threat from this new malware, we also discuss possible countermeasures.

5.1. Extension of the Malware

We have shown that malware activation (the application launch detection) rate is 100%, and any variety of malicious payloads may be triggered this way. This opens the door for attacks with an increased amount of sophistication, an alarming thought given Zhou and Jiang's [33] evaluation of various state of the art mobile anti-virus applications.

The proof of concept malware described in this article can be improved in a number of ways. For example, our current Activity does nothing to hide itself from the Recent Tasks list, the window which displays a list of the most recently used applications. To become much harder to detect, a hacker could supply the `FLAG_ACTIVITY_EXCLUDE_FROM_RECENTS` flag in the application manifest to exclude the malicious Activity from the Recent Tasks list. Another area of improvement is the type of events we are processing. Our example focuses on AccessibilityEvents of type `TYPE_VIEW_CLICKED`, however there are 22 other types that could potentially be used to activate malicious payloads with 100% success rate. We could improve the likelihood of activation by matching on more application launches. In fact, our technique could be used to gather information about packages installed

on the device without the need to request permission from the Package Manager. It could do so by harvesting data from the package, class, and text fields of the AccessibilityEvents it receives. Simply collecting the text field of the AccessibilityEvent, as we do during launch detection, will provide information about which apps are installed on the device. Combined with the remote payload update technique, this malware could be dynamically extended to provide a malicious payload for every app installed on the device. Finally, this attack could be used for denial of service or ransomware. An attacker could detect the launch of the Settings application and display a blank Activity, making the user unable to change any device settings. Moreover, an attacker could detect the launch of any application (even the Launcher itself) and display their own ransom Activity, effectively rendering the device useless until the attacker is paid some amount of money.

5.2. Countermeasures

The attacks mentioned above are not foolproof. There are certain safeguards built into Android devices to thwart a full device takeover by a malicious user. Recovery Mode is one such safeguard that allows the user to install a clean OS, wiping any malicious apps from the data partition in the process. An experienced user could use the Android Debug Bridge (ADB) to obtain a command shell into the device and launch applications from the command line *am* tool or even locate and uninstall the malicious applications. Applications launched from the *am* tool do not invoke the Launcher application, therefore our technique is unable to detect this. If malicious application developers were to become aware of the *am* countermeasure, the malicious AccessibilityService could be rewritten to trigger on a non-Launcher event. For example, any Activity containing a button widget that launches a new Activity could be targeted if the malware is modified to trigger on the button click AccessibilityEvent rather than the Launcher event. However, doing so may limit the generic nature of this attack.

As a countermeasure to the malicious Activity payload, the logic which reorders the Activity history stack and determines the next Activity to be displayed should be analyzed and corrected. The two main Android files containing said logic are `frameworks/base/services/java/com/android/server/am/ActivityManagerService.java` and `frameworks/base/services/java/com/android/server/am/ActivityStack.java`. Ensuring that the next Activity to be displayed is indeed the Activity that was requested by `startActivity` could prevent the malicious Activity from being displayed instead of the legitimate Activity. However, this fix does not hinder launch detection and the attacker could

simply delay the start of the malicious Activity for some short time, ensuring that the malicious Activity is displayed after all. The difference here is that there may be some obvious transition animation, if the developer of the legitimate app has not disabled it, from the legitimate Activity to the malicious Activity.

6. Related Work

In 2009, Schmidt *et al.* [31] made a survey of mobile malware and found that most malware targets Symbian OS. They reported F-Secure Research in Helsinki counted 418 malware samples, some of which were not public, while they collected information of 288 public malware. 278 of these 288 public malware targeted Symbian OS. A note is On February 11 2011, Nokia announced to adopt Microsoft's Windows Phone OS as its primary smartphone platform, and Symbian has faded out since then. Since Android OS was getting attention at that point and the authors investigated possibilities of malware on Android, they explored "social engineering"-based Android malware, where the malicious functionality is hidden in a seemingly benign host app. They demonstrated such functionality can be binary code. *android.os.Exec* can be used to finally execute such binary code. The binary code is the payload of the malware. The authors show the payload can be crafted to bypass the Android permission system such as accessing */proc* and */sys/*, deplete the device's battery by using energy consuming FPU (Floating Point Unit) operations, and run arbitrary ARM instructions on a rooted G1 Android smartphone. Felt *et al.* [26] classify threats from third-party smartphone applications into *malware*, *grayware*, and *personal spyware*. Malware intends to damage finance or property of the smartphone owner. An "attacker" such as a spouse who has physical access to a smartphone can install personal spyware on the victim smartphone and gather information about the smartphone owner, for example, tracking the victim. Grayware is often commercial applications with real functionality while stealing user information. The distributor may have a privacy policy with varying degree of clarity. The authors conduct a survey of 46 pieces of smartphone malware and their incentives and conclude that Apple's mechanisms of application permission and review process can avoid approving malware. Becher *et al.* [22] examine mechanisms securing sophisticated mobile devices in 2011. Although no major incidents of attacking smartphones have happened, small-scale attacks have been emerging. Threats were classified into four classes: *hardware centric*, *device independent*, *software centric*, and *user layer attacks* for the purpose of eavesdropping, availability attacks, privacy attacks and impersonation attacks. Existing security mechanisms are enumerated for various attacks.

Enck *et al.* [25] implemented *ded*, a Dalvik decompiler. *ded* transfers *.dex* file into Java source code. It was then used for analyzing security of 1,100 popular free Android applications. The following major observations were made: misuse of privacy sensitive information including phone identifiers such as IMEI, IMSI, and ICC-ID and geographic location; "no evidence of telephony misuse, background recording of audio or video, abusive connections, or harvesting lists of installed applications"; wide use of ad and analytic network libraries by 51% of the applications; no exploitable vulnerabilities leading to control of the phone. Zheng *et al.* [32] developed ADAM, an automated system for evaluating the detection of Android malware. ADAM uses repackaging and code obfuscation to generate different variants of a malware. They collected 222 malware samples and used ADAM to generate variants of those malware. Those variants were fed into VirusTotal [21], "a free service that analyzes suspicious files and URLs and facilitates the quick detection of viruses, worms, trojans, and all kinds of malware". They have observed that commercial anti-virus showed different detection rate for different variants. New anti-virus software such as Antiy [14] shows better performance than older anti-virus software. Rastogi, Chen and Jiang [30] made similar effort to test state-of-the-art Android commercial mobile anti-malware products for detecting transformed malware. Such transformation techniques include polymorphism (where transformed code is still similar to the original code) and metamorphism (where transformed code is totally different from the original code, but with similar malware functionality). Zhou *et al.* [34] developed a system called DroidRanger, evaluating the health of Android markets, including the official Android Market, eoe-Market [16], alcatelclub [12], gfan [17], and mmoovv [20]¹. To detect known Android malware, DroidRanger uses permission-based filtering to detect malware using suspicious permissions and behavioral footprint matching to detect malware performing suspicious behavior such as listening to system-wide broadcast messages and sending and monitoring SMS messages. To detect unknown Android malware, DroidRanger uses two steps: heuristics based filtering and dynamic execution monitoring. The heuristics based filtering can utilize Android features misused to load new code, either Java binary code from a remote server or native machine code. Dynamic execution monitoring checks what APIs an app is using. DroidRanger was able to find 211 malicious or infected apps out 204,040 apps from the five studied marketplaces, including two zero-day malware. Zhou and Jiang [33] made a one year effort and analyzed more than 1,200 malware samples,

¹Link is no longer valid

which covered a majority of state-of-the-art Android malware. They obtained those samples by manually or automatically crawling various Android markets. They characterize Android malware by their installation, activation, payload and permission use. Installation of malware on a victim device uses three main social engineering based approaches, repackaging, update attack, and drive-by download. In repackaging, a benign app is downloaded, piggybacked with malicious code and uploaded onto a market again. In update attack, a malware author put code into an app, and the rest of the malicious code will be downloaded when the malware is running. In drive-by download, ad is used in a malware to attract the victim to download more spyware and other malware. Other attacks also exist: spyware, fake apps masquerading as other legitimate apps, apps with malicious functionality such as sending unauthorized SMS messages, apps exploiting root privilege. Malicious apps can be activated by various system events including `BOOT_COMPLETED`, `SMS_RECEIVED`, and UI interaction events. Malware can have a variety of payloads, targeting privilege escalation, remote control, financial charges, and personal information stealing. Malware without root exploits often utilizes `INTERNET`, `READ_PHONE_STATE`, `ACCESS_NETWORK_STATE`, and `WRITE_EXTERNAL_STORAGE` permissions. The authors have found that malware have been evolving to avoid detection and have more sophisticated functionality such as making the device part of a botnet. Current anti-virus software downloadable from Google market, AVG Antivirus Free, Lookout Security & Antivirus, Norton Mobile Security Lite, and Trend Micro Mobile Security Personal Edition, do not perform well in detecting malware the authors collected. Bugiel *et al.* [24] studied ways to defend against privilege-escalation attacks on Android. Such privilege-escalation attacks include confused deputy attacks and colluding attacks. Confused deputy attacks exploit unprotected interfaces of a benign application. Colluding attacks involve multiple apps. For example, one app can record audio and another one has the Internet permission. In this way, the second app can send the overheard credit numbers out. The authors designed and implemented a security

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