

Advanced Data Structures

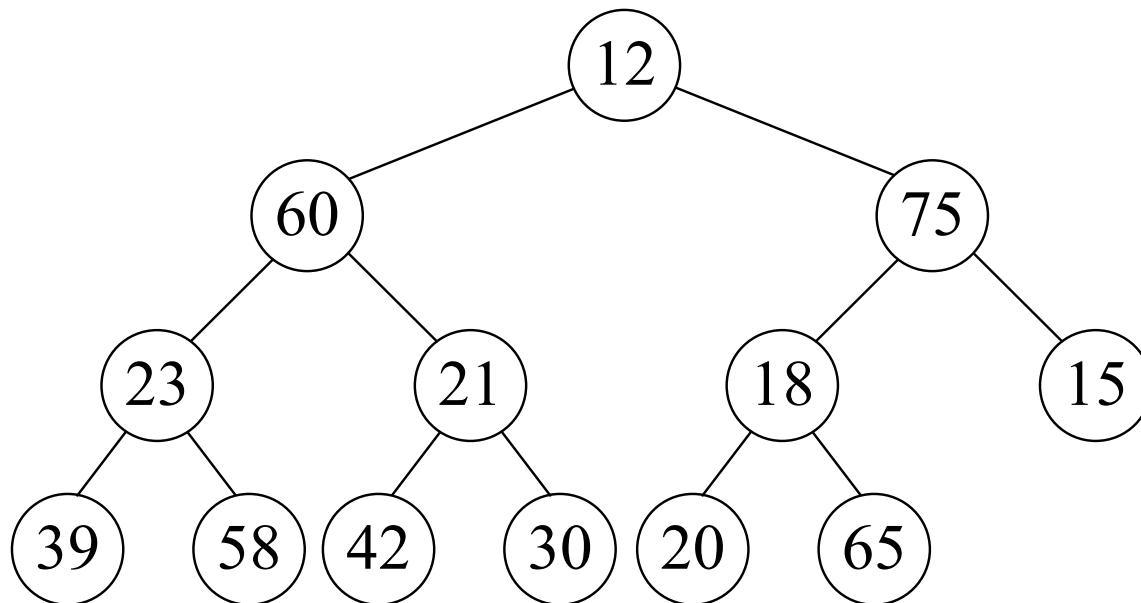
Min-Max Heap

Min-Max Heap

23



■ l e 10 80



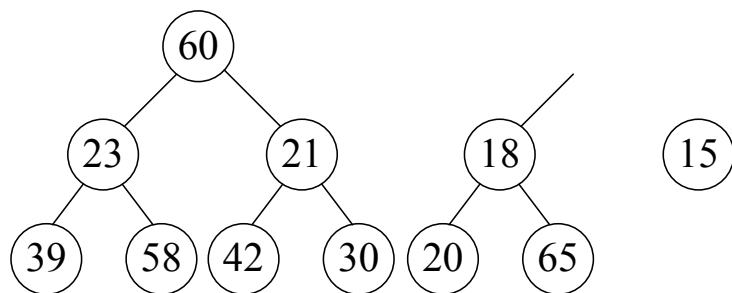
min level

ma level

min level

ma level

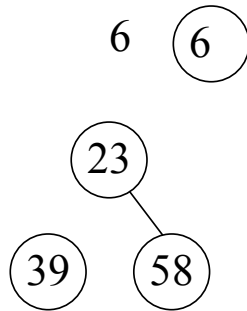
1.



2.

3.

4.



5. Ne de i i 9 na le el a d i la ge ha
i a e 9 e cha ge

6.

```

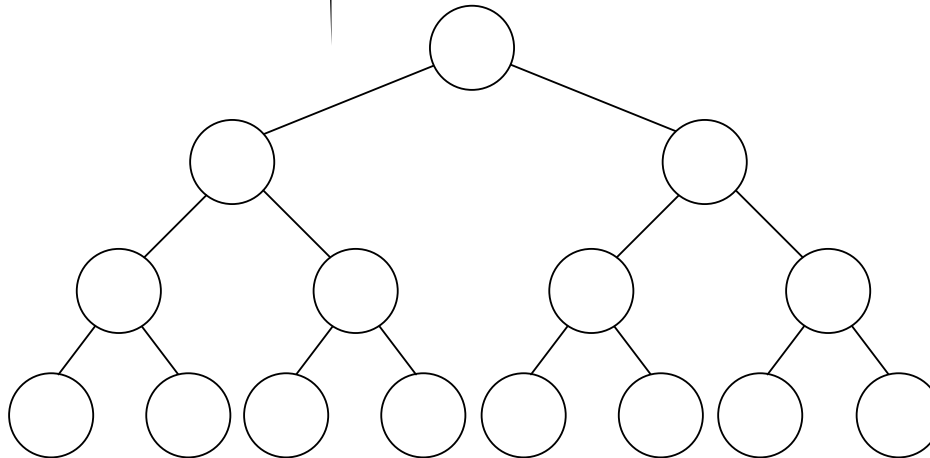
template <class T pe>
void MinMa_Heap<T pe>::Insert (const Elemen<T pe>& e)
    if (n == Ma_Size) MinMa_Fill(); return;
    n++;
    int p = n/2;
    if (!p) h[1] = e; return;
    switch (level(p))
    case MIN:
        if (e.ke < h[p].ke)
            h[n]=h[p];
            Verif_Min(p, e);

```

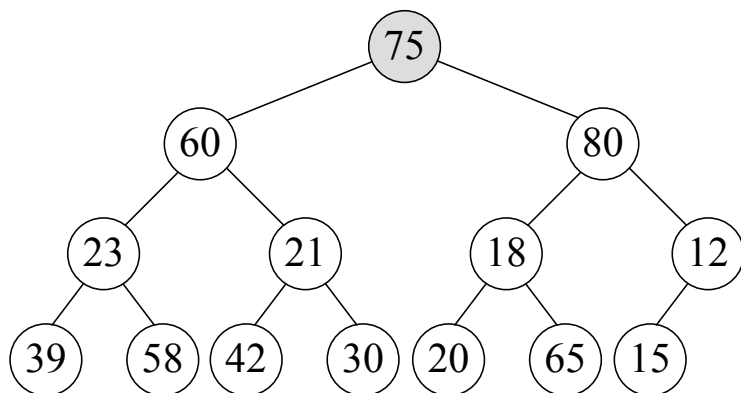
```
    else Verif Ma (n, );  
    break;  
case MAX:  
    if ( .ke > h[p].ke )  
        h[n]=h[p];  
        Verif Ma (p, );  
  
    else Verif Min(n, );
```

- **Level:**
Ma Level:
 $\lceil \log_2(j + 1) \rceil$ is even
Min Level:
 $\lceil \log_2(j + 1) \rceil$ is odd
- **Verif Ma**
- **Verif Min**

- Delete Min from Min-Heap



1.



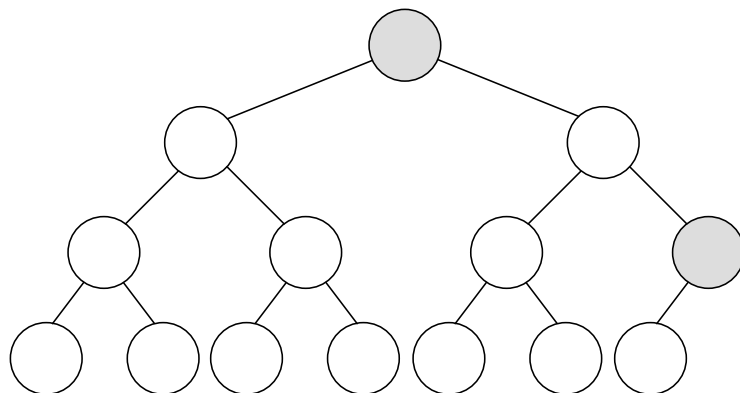
min level

max level

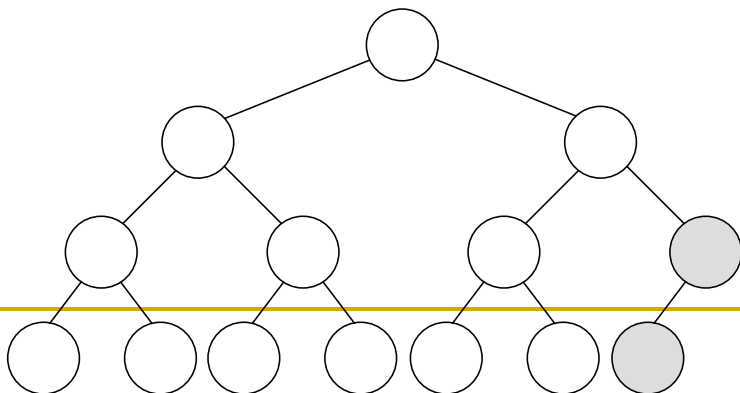
min level

max level

2.



3.



I e a e de i a i - a hea i h

1 e hea : i he e ;

2 a lea e child f :

fi d de k i h he alle da a;

(a) .ke h[k].ke

i he e ;

(b) .ke > h[k].ke a d k i a child f
:

 e lace k i h ;
 k i he e .

(c) .ke > h[k].ke a d k i a g a dchild
f :

 k i he e

 if .ke > h[].ke e cha ge a d h[]

 i e i (b) i _ a hea ed
i h k

```
template <class T pe>
```

```
Elemen <T pe>* MinMa Heap<T pe>::
```

```
    Dele eMin(Elemen <T pe>& )
```

```
    if (!n) MinMa Emp ( ); return 0;
```

```
        = h[1];
```

```
    Elemen <T pe>    = h[n--];
```

```
    int i = 1, j = n/2;
```

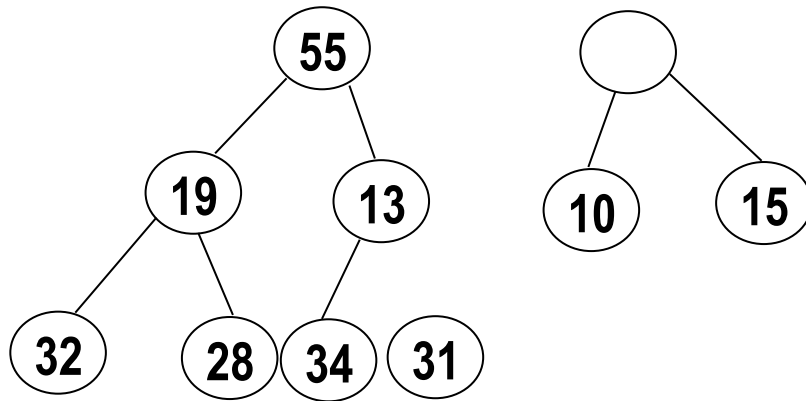
```
        hile (i <= j)    // ha children
```

```
        int k = MinChildGrandChild(i);
```

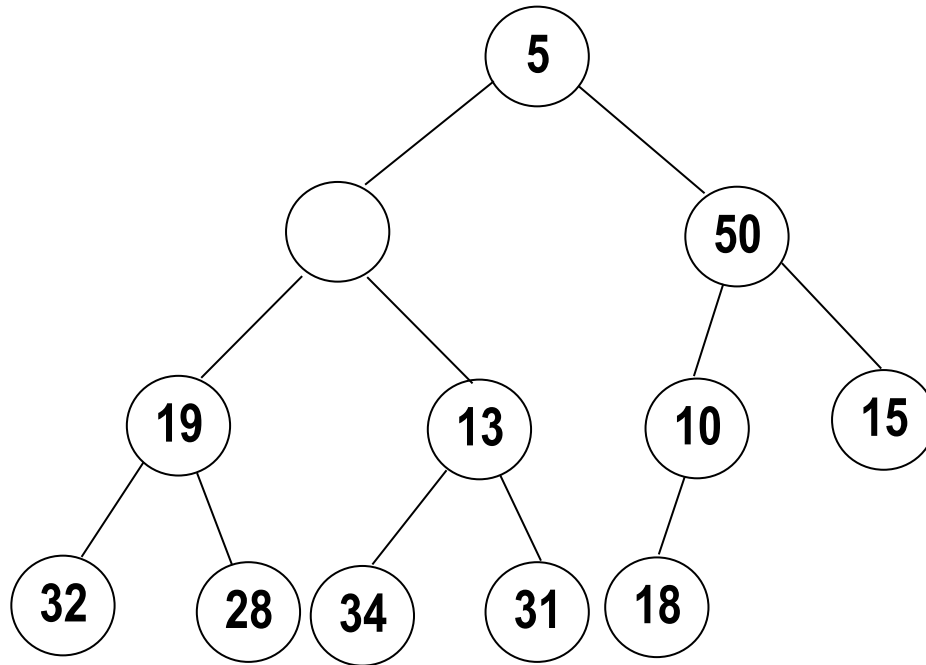
```
        if ( .ke <= h[k].ke )
```

```
            break;    // 2(a)
```


Min-max heap



Min-max heap (con.t)



❑ Dele e 55

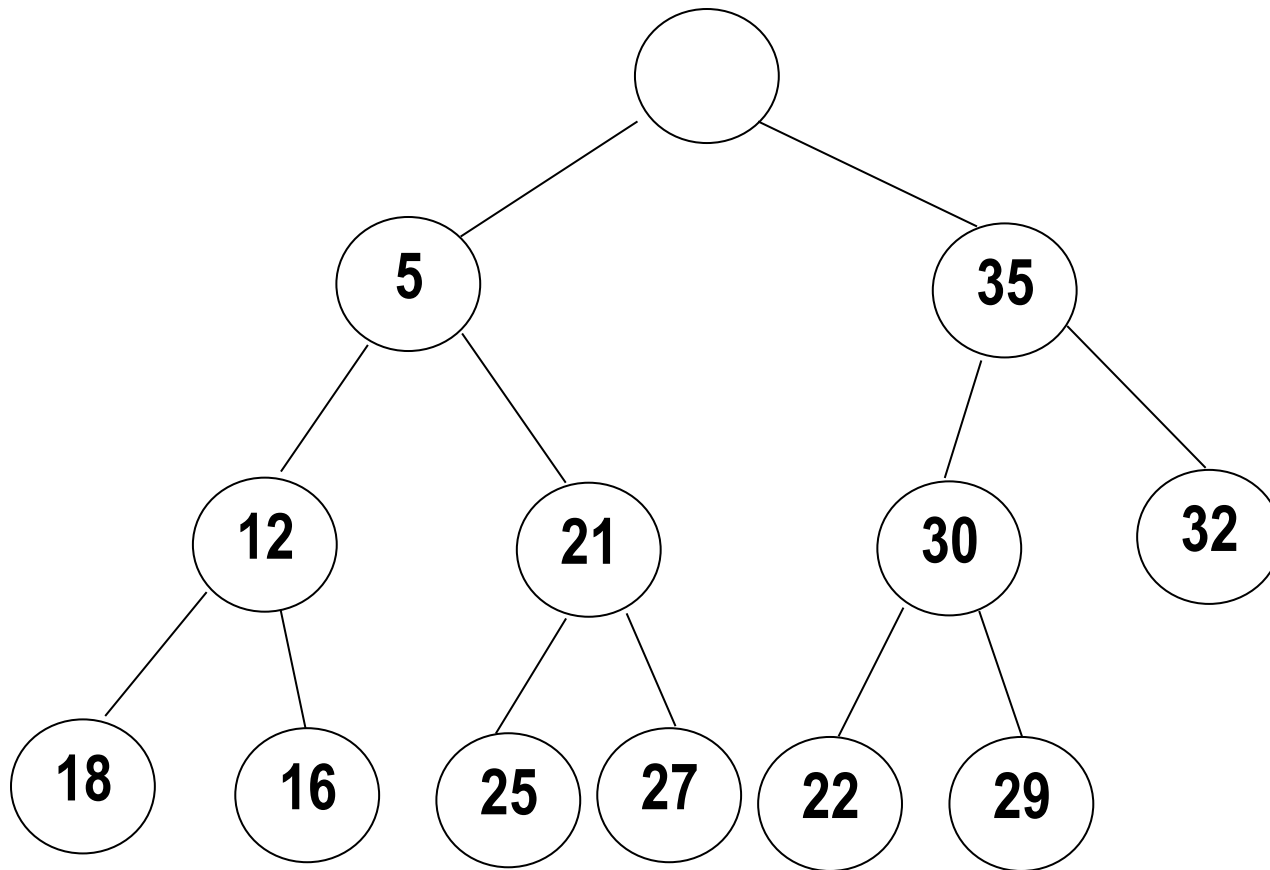
Advanced Data Structures

Deap

A deap is a complete binar tree:

- (1) no data in root;**
 - (2) left subtree is a min heap**
 - (3) right subtree is a ma heap**
 - (4) ever node in left subtree is no bigger than its
corresponding node in right subtree
corresponding nodes**
-

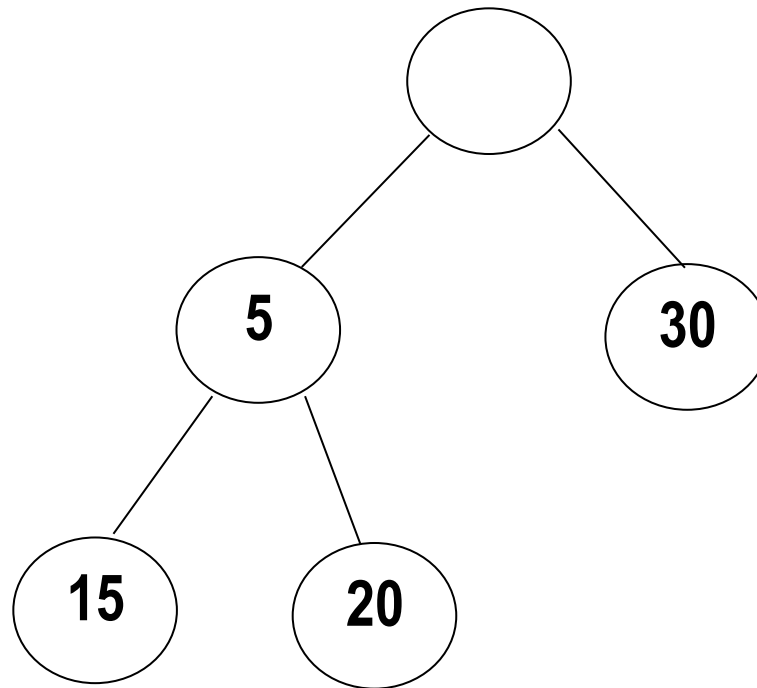
Deap



Deap (con.t)

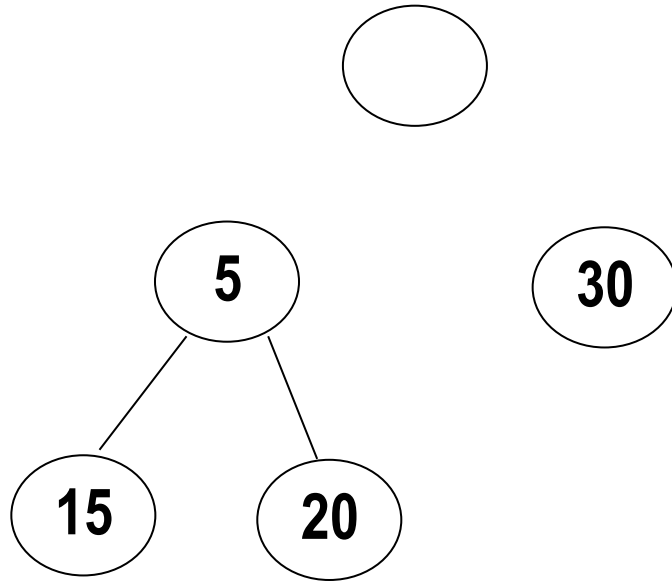
■ l e i

□ Add e de he e d f ee, **Heapif**

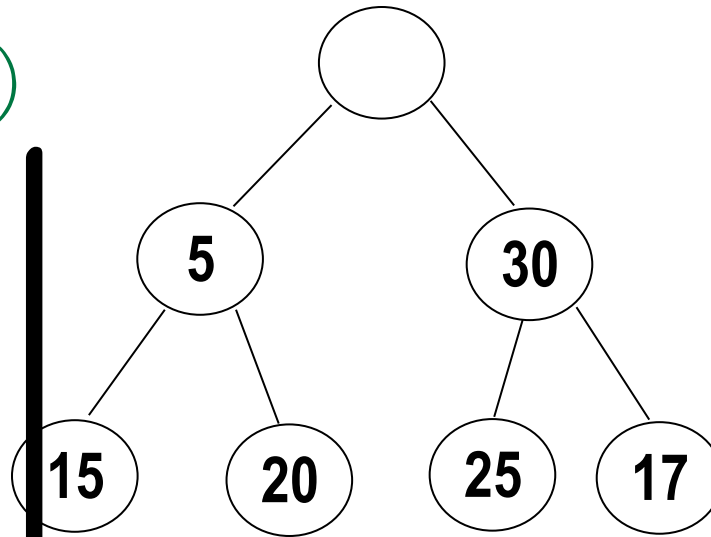


□ l e 25

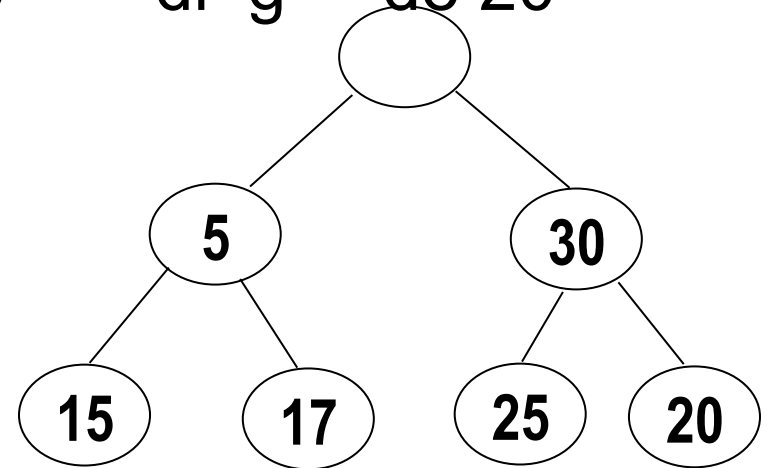
Deap (con.t)



Deap (con.t)

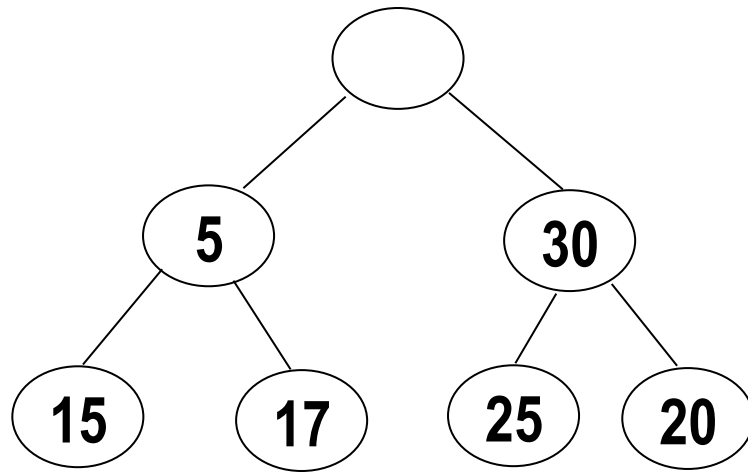


- 17 is smaller than its parent 30, so it is a local minimum.
- Exchange 17 and 20.
- Check the new heap.



- Need to check the new heap, WHY?

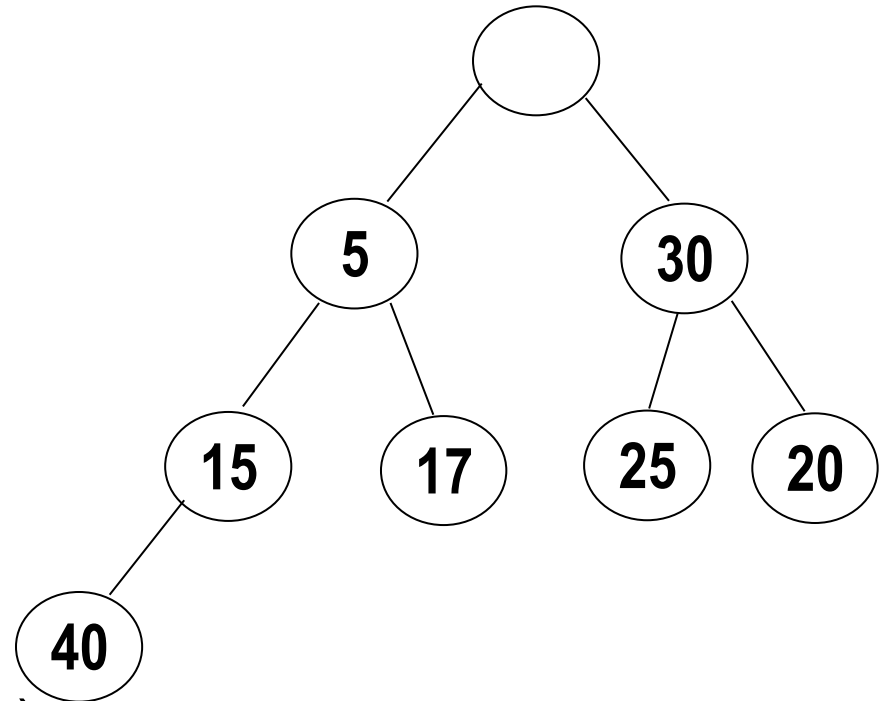
Deap (con.t)



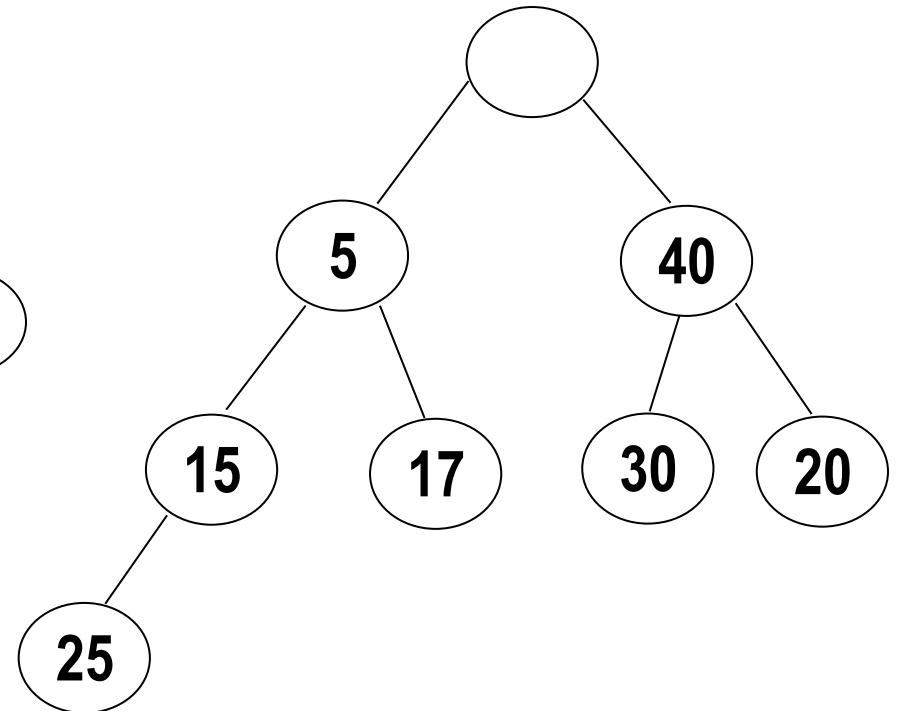
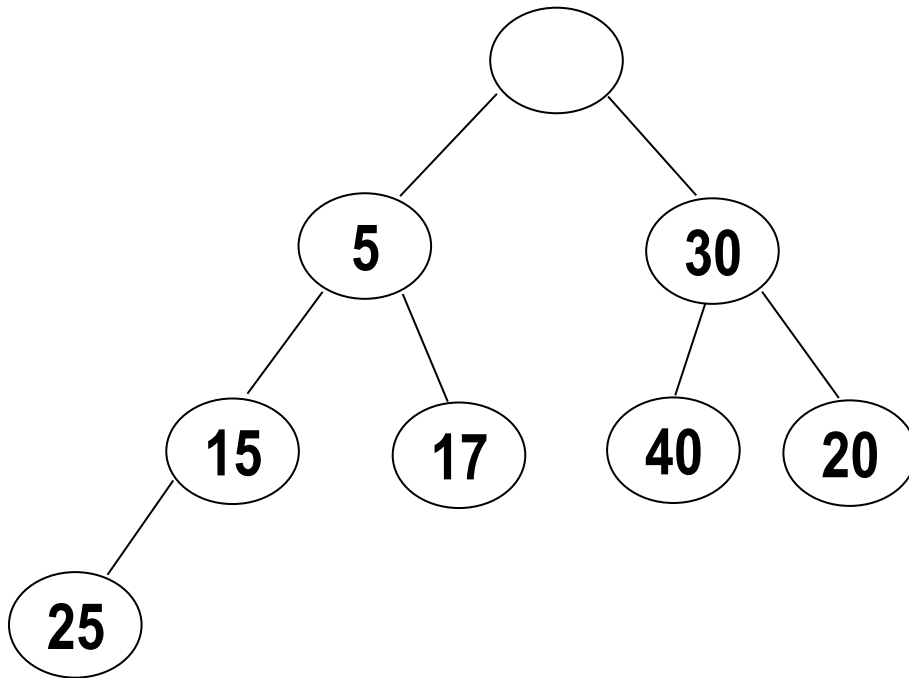
□ I e 40

□ 40 i bigge ha 25()

9 E cha ge 40 a d 25



Deap (con.t)



❑ The right subtree is not a max-heap

■ **Heapify**

❑ No need to adjust the left subtree, WHY ?

```

template <class T pe>
void Deap<T pe>::In er (const Elemen <T pe>& )

    int i;
    if (n == Ma Si e ) DeapF ll( ); return;
    n++;
    if (n == 1) d[2] = ; return;
    int p = n + 1;
    switch (Ma Heap(p))
        case TRUE: // p in ma -heap
            i = MinPar ner(p);
            if ( .ke < d[i].ke )
                d[p] = d[i];
            MinIn er (i, );

```



(1) Deap::Ma Heap(int p)
for $p > 1$, if $2^{\lfloor \log_2 p \rfloor} + 2^{\lfloor \log_2 p \rfloor - 1} \leq p < 2^{\lceil \log_2 p \rceil}$, p
is in the ma -heap

(2) Deap::MinPartner(int p)

P is in ma -heap and P is the last one in Deap:

$$p - 2^{\lfloor \log_2 p \rfloor - 1}$$

(3) Deap::Ma Partner(int p)

P in min-heap and P is the last one in Deap:

$$(p + 2^{\lfloor \log_2 p \rfloor - 1}) / 2$$

